Tight Bounds on Minimax Regret under Logarithmic Loss via Self-Concordance

Blair Bilodeau^{1,2,3}, Dylan J. Foster⁴, and Daniel M. Roy^{1,2,3} Presented at the 2020 International Conference on Machine Learning

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Contextual Online Learning with Log Loss

Example: Image Identification

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Challenges

- We do not rely on data-generating assumptions.
- ℓ_{\log} is neither bounded nor Lipschitz.

Measuring Performance with Regret

Without model assumptions, guaranteed small loss on predictions is impossible.

Consider a relative notion of performance in hindsight.

- Relative to a class $\mathcal{F} \subseteq \{f : \mathcal{X} \to [0,1]\}$, consisting of **experts** $f \in \mathcal{F}$.
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Regret:
$$R_n(\hat{\mathbf{p}}; \mathcal{F}, \mathbf{x}, \mathbf{y}) = \sum_{t=1}^n \ell_{\log}(\hat{p}_t, y_t) - \inf_{f \in \mathcal{F}} \sum_{t=1}^n \ell_{\log}(f(x_t), y_t).$$

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This quantity depends on

- p: Player predictions,
- \mathcal{F} : Expert class,
- x: Observed contexts,
- y: Observed data points.

Summary of Results

We control the **minimax regret** using the **sequential entropy** of the experts \mathcal{F} .

- Minimax regret: the smallest possible regret under worst-case observations.
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Contributions

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- Improved upper bound for expert classes with polynomial sequential entropy.
- Novel proof technique that exploits the curvature of log loss to avoid a key "truncation step" used by previous works.
- Resolve the minimax regret with log loss for Lipschitz experts on $[0,1]^p$ with matching lower bounds.
- Conclude the minimax regret with log loss cannot be completely characterized using sequential entropy.

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The first context is observed.

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The player makes their prediction.

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The adversary plays an observation.

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Interpretation: The experts \mathcal{F} are *minimax online learnable* if $R_n(\mathcal{F}) < o(n)$.

- slow rate: $R_n(\mathcal{F}) = \Theta(\sqrt{n})$
- fast rate: $R_n(\mathcal{F}) \leq \mathcal{O}(\log(n))$

Covering Numbers

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Uniform Covering

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A uniform covering may be infinite for large expert classes. Instead, we use *sequential covering* from Rakhlin and Sridharan (2014).

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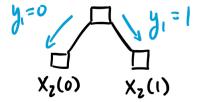
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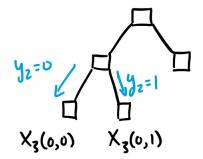
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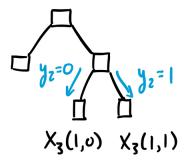
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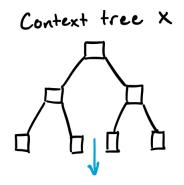
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A class of trees V sequentially covers ${\mathcal F}$ at margin γ on context tree ${\bf x}$ if:

$$\sup_{f \in \mathcal{F}} \sup_{\mathbf{y} \in \{0,1\}^n} \inf_{\mathbf{y} \in V} \sup_{t \in [n]} |f(x_t(\mathbf{y})) - v_t(\mathbf{y})| \le \gamma.$$

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Definitions

- The size of the smallest such V for x is $\mathcal{N}_{\infty} (\mathcal{F} \circ \mathbf{x}, \gamma)$.
- Sequential entropy for n rounds is $\mathcal{H}_{\infty}(\mathcal{F}, \gamma, n) = \sup_{\mathbf{x}} \log (\mathcal{N}_{\infty}(\mathcal{F} \circ \mathbf{x}, \gamma)).$

Theorem (BFR '20)

There exists c > 0 such that for all \mathcal{F} ,

$$R_n(\mathcal{F}) \leq \inf_{\gamma > 0} \Big\{ 4n\gamma + c \,\mathcal{H}_{\infty}\left(\mathcal{F}, \gamma, n\right) \Big\}.$$

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Theorem (BFR '20)

If $p \in \mathbb{N}$, there exists an \mathcal{F} with $\mathcal{H}_{\infty}(\mathcal{F}, \gamma, n) = \Theta(\gamma^{-p})$ and

 $R_n(\mathcal{F}) \ge \Omega(n^{\frac{p}{p+1}}).$

• 1-Lipschitz:

 $\mathcal{F} = \{ f \mid f : [0,1]^p \to [0,1], |f(x) - f(y)| \le ||x - y|| \quad \forall x, y \in [0,1]^p \}.$ $\mathcal{H}_{\infty} \left(\mathcal{F}, \gamma, n\right) = \Theta \left(\gamma^{-p}\right).$

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We have matching upper and lower bounds for this class, so:

$$R_n(\mathcal{F}) = \Theta(n^{\frac{p}{p+1}}).$$

• Linear Predictors:

$$\mathcal{F} = \{ f \mid \exists w \text{ s.t. } \|w\|_2 \le 1, f(x) = \frac{1}{2} [1 + \langle w, x \rangle] \; \forall \, \|x\|_2 \le 1 \}.$$

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Our upper bound cannot be improved, so the minimax regret under log loss cannot be characterized solely by sequential entropy.